



# ELF II NEWSLETTER

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### The Great Micro Power Contest by Tom Pittman

Most of you know by now that the February issue of Kilobaud has an article by me entitled "DOTS". In this article I explain how the character generator in Netronics Tiny Basic works. At the end of the article I put out a challenge: "The 1802 is the most powerful 8-bit microprocessor," I said, and it was worth \$50 out of my pocket to prove it. Well, I was kind of dumb in laying out the rules for the prize.

Basically it was this: DOTs is a callable subroutine to display characters on the 1861, one character per call. I got it into some 800-odd bytes of code, with an average execution time around 1300 or so memory cycles (i.e. something less than 700 instructions). Multiply these together for a processor power rating of about 1.1 million byte-cycles, or "bykes" as I will call them here. The rules for winning the \$50: beat that by 10%.

Well, somebody did. Before I even got my copy of the February Kilobaud, I received a program from someone in Chicago (he said he was stuck in a blizzard) which was written for a Z80 and came to 32% less bykes than mine. I paid him his \$50, but not before I looked at his program veeerrrry carefully. You see, he changed the data structure of the program, so he did not prove the Z80 more powerful than the 1802. What he proved is that Mike Amling is a more clever programmer than Tom Pittman.

How do I know? I reworked my program along the same lines and beat his program by another 0.2%. OK, obviously the Z80 is as good as the 1802. But not 30% better.

The reason I mention all this is that all of you who are using the 1802 will be subjected to various forms of bigotry and prejudice, by people who will try to tell you that the 1802 is less powerful than Brand X. Don't you believe it. They just made the mistake of buying into a less powerful CPU and are trying to cover their tails.

When you get right down to it, the differences are slight. The average instruction length in the 1802 is about 1.3 bytes. In the 8080 it is closer to 1.9 bytes. In the 6800 I would guess 1.7 bytes in a typical mix of instructions. That makes the 1802 look a lot better, but you have to realize that because the instructions are smaller it takes some 30% more of them to do the same job, bringing the effective value to about 1.7. Notice that this is still very competitive.

Now look at execution time. But to be fair, let's not measure an 1802 with a 1.7MHz clock against a 6502 with a 4MHz clock. When it was first announced, the 1802 was (and still is) specified at 6.4MHz, for a memory cycle time of 1250ns. The 8080 was specified for a clock speed of 2MHz, giving a memory cycle time of 1500ns. The original Z80s were specified for the same speed, though very shortly they were announcing 4MHz versions. The original 6800 and 6502 spec called for a 1MHz

clock and 1000ns memory cycle (these also have speeded up). The RCA people do not seem to understand what sells, so they have not bothered to make a speedup version of the 1802. More's the pity.

But let's consider the original releases of each of these. On the 1802 no instruction is more than three memory cycles, and most are exactly two cycles. That's 2.5 microseconds on a plain-vanilla 1802 running maximum clock. On the 8080 and Z80 there are some instructions that execute in four clock cycles (2us) but most of them require 5, 7, 8, or often 11 and 17 clock cycles to execute. The most common instructions (moves and conditional jumps) generally run four and five microseconds each. In the 6800 the figures are very similar: some instructions execute in only 2 memory cycles; most take three to five cycles, and one takes twelve.

The Z80 has some sexy instructions like block move and block I/O. A two-byte instruction can move any number of bytes at the rate of one byte every 21 clock cycles (10.5us in the plain-vanilla model). The 1802 can do the same thing in a 7-byte loop in 15us; if I get to choose where the data is coming from or going to, the loop is six bytes and only 12.5us per byte moved. The block I/O instructions take the same time in the Z80; in the 1802 I can code a four-byte loop to do the same thing in as little as 7.5 microseconds per byte of output.

Well, then, why isn't the 1802 more popular? "It's the addressing modes," say the armchair critics. "The 6502 has 13 addressing modes and the 1802 has hardly any." Six modes is not exactly "hardly any" (the 8080 has only five). But I can program any addressing mode I want into the 1802. The ELFBUG program uses a (programmed) relative addressing mode. I usually program a base-page addressing mode into most of my programs. You could write 1802 programs with 23 addressing modes if you wanted (and if you could think of a use for them); in the 6502 you are stuck with the 13 they give you.

"But the instruction set is so ridiculous," the critics splutter. I say, who is ridiculing whom? The 1802 can do a subroutine call in 2.5 microseconds; no other micro does it in less than five. The 1802 can receive an interrupt, process it, and be back in the main program in the time it might take the 6800 to notice it got an interrupt. Some people insist that the 6502 "is more like a minicomputer such as the PDP-11." Izzatso? Show me a 6502 program that can do arithmetic on the program counter. Only the PDP-11 and the 1802 can do that. I claim the 1802 also has many of the important features found only in large computers like the DEC-10 and the 360.

I have only begun to list the advantages of the 1802, but I think I had better quit before I get too excited. I just wanted you to know that you have nothing to be ashamed of. Now if only we could convince the folks at RCA to give the 1802 the support it deserves.

## The Elusive 256

The basic ELF II comes with 256 bytes of RAM. This is enough to write a large variety of programs, but not quite enough to run bigger things like Tiny Basic. Enter 4K. For a very reasonable price, you can add 4K RAM in the bus connector. Tiny will run OK in 4K, but it does not leave you much room for programs. However, if you leave the 256 bytes in place, you get 4.3K, which is a lot more program space. The problem is, it did not work all the time. I fixed mine by changing two resistors on the memory board: Change R6 from 10K to 1.2K on every memory board; and add R5 (this was left out of my boards) 3.9K on one memory board only. Can you figure out why this works? One of these days I will have to do an article on hardware debugging. Promises, promises.

### Multiple Timer Program

by J. H. Hansen

This program was written for use in the photographic darkroom where multiple times of different lengths are required.

Features of the program are:

1. As written the program provides for eight different times. This could be expanded to 18+ if all memory locations are used.
2. The memory locations for desired times are near the beginning of the program, making it very easy to change them.
3. The times are entered and displayed in decimal form, that is, directly in minutes and seconds. Therefore no knowledge of or conversion to hexadecimal is required.
4. Remaining minutes are displayed and when they are zero, remaining seconds are displayed. When remaining seconds are 15 seconds or less a warning beep is sounded each second. At End of Time a tone of different pitch and longer duration is sounded.
5. To start the timer it is only necessary to push the "I" button twice. When time is finished, the next time will start when the "I" button is again pushed twice. The series of times can be reset to the beginning by turning the Run switch off and on.
6. The maximum time limitations are 99 minutes and 99 seconds. If you are willing to work with hexadecimal numbers, the program could be made shorter and the time limits expanded considerably. [But who would want to? —TP]

[Editor's note: This program requires a speaker to produce the sounds. An audible sound can be produced by connecting a cheap speaker across the Q LED. A much louder sound may require an amplifier.]

Addr. Hex	Instruction	Comments
00 F800	LDI 00	Initialize high order
02 B4	PHI R4	registers 4,5,6,7.
03 B5	PHI R5	
04 B6	PHI R6	
05 B7	PHI R7	
06 F826	LDI MAIN	Set memory location
08 A4	PLO R4	for main program
09 F8AF	LDI DECH	Set memory location for
0B A5	PLO R5	Decimal to Hex subroutine
0C F8C6	LDI HEXD	Set memory location for
0E A6	PLO R6	Hex to Dec subroutine

0F F8FF	LDI #FF	Set memory location for
11 A7	PLO R7	decimal data storage
12 D4	SEP R4	Call main program
13 MM SS		Enter times desired
15 MM SS		Minutes & seconds in Decimal
17 MM SS		max 99 Min. & 99 Sec.
19 MM SS		Eight different times
1B MM SS		may be entered
1D MM SS		
1F MM SS		
21 MM SS		
23 F810	NEXT: LDI #10	Set tone for
25 5D	STR RD	warning "beep"
26 E0	MAIN: SEX R0	
27 3F27	BN4 *	Press Input button
29 F0	LDX	Get minutes
2A 64	OUT 4	Display minutes
2B E7	SEX R7	
2C D5	SEP R5	Convert decimal to hex
2D B3	PHI R3	Put minutes in Reg 3
2E 372E	B4 *	Release button
30 3F30	BN4 *	Push Button
32 E0	SEX R0	
33 F0	LDX	Get seconds
34 64	OUT 4	Display seconds
35 E7	SEX R7	
36 D5	SEP R5	Convert decimal to hex
37 A3	PLO R3	Put seconds in Reg 3.0
38 3738	B4 *	release button
3A 7B	SEQ	Turn "Q" LED on
	.. One minute loop. Repeat until Min=0	
3B 93	MIN: GHI R3	Get minutes
3C D6	SEP R6	Convert hex to decimal
3D 57	STR R7	Put min in loc.FF
3E 64	OUT 4	Display minutes
3F 27	DEC R7	
40 93	GH1 R3	Get min.
41 3259	BZ SEC	If Min=0
43 FF01	SWI 01	
45 B3	PHI R3	
46 F83C	LDI 60	
48 AE	PLO RE	
49 F88D	LDI (37261).0	
4B AF	ML: PLO RF	
4C F892	LDI (37261).1	
4E BF	PHI RF	
4F 2F	DEC RF	
50 9F	GH1 RF	
51 3A4F	BNZ *-2	
53 2E	DEC RE	
54 8E	GLO RE	
55 3A49	BNZ ML	
57 303B	BR MIN	
	.. 1 second loop. Repeat until sec.<15	
59 83	SEC: GLO R3	GET SECONDS
5A D6	SEP R6	Convert hex to dec.
5B 57	STR R7	Put sec. in loc.FF
5C 64	OUT 4	Display seconds
5D 27	DEC R7	

```

5E 83      GLO R3
5F 32A1    BZ OUT If Sec=0
61 FF0F    SMI 15 Subtract 15 from D (sec)
63 3B72    BNF BEEP If 0 or less
65 23      DEC R3
66 F8A7    LD1 (37287).0
68 AF      PLO RF
69 F892    LD1 (37287).1
6B BF      PHI RF
6C 2F      DEC RF
6D 9F      GHI RF
6E 3A6C    BNZ *-2
70 3059    BR SEC

.. Part of Second timing loop with warning beep
.. every second for last 15 seconds.

72 83      BEEP:GLO R3 Get sec.
73 D6      SEP R6 convert to sec.
74 57      STR R7 Store in memory loc.FF
75 64      OUT 4 Display sec.
76 27      DEC R7
77 7A      REQ
78 F803    TEND:LD1 (521).1
7A BC      PHI RC
7B F809    LD1 (521).0
7D AC      PLO RC
7E F810    SL: LD1 #10
80 A8      PLO R8
81 7B      SEQ
82 FF01    SMI 01
84 3A82    BNZ *-2
86 7A      REQ
87 88      GLO R8
88 FF01    SMI 01
8A 3A88    BNZ *-2
8C 2C      DEC RC
8D 9C      GHI RC
8E 3A7E    BNZ SL
90 83      GLO R3
91 3223    BZ NEXT If Sec=0 wait for next time

.. Part of Second timing loop
.. to make loop = 1 second

93 F851    LD1 (25169).0
95 AF      PLO RF
96 F863    LD1 (25169).1
98 BF      PHI RF
99 2F      DEC RF
9A 9F      GHI RF
9B 3A99    BNZ *-2
9D 23      DEC R3
9E 83      GLO R3
9F 3A72    BNZ BEEP
A1 57      OUT: STR R7
A2 64      OUT 4
A3 27      DEC R7
A4 F87F    LD1 SL+1 Location of tone for
A6 AD      PLO RD for signalling end of time
A7 90      GHI R0
A8 BD      PHI RD
A9 F850    LD1 #50 End of time note (tone)

```

```

AB 5D      STR RD Go back thru second loop
AC 3078    BR TEND to signal end of time

.. Decimal to Hex Subroutine
AE D4      SEP R4 Go back to main program
AF AA      DECH:PLO RA Enter byte
B0 FAF0    ANI #F0 Find most sig. Hex digit
B2 F6F6F6  SHR 3 MSHD x2
B5 73      STXD
B6 60      IRX
B7 F4F4F4F4 ADD;ADD;ADD;ADD
BB 73      STXD
BC 60      IRX
BD 8A      GLO RA Get original byte
BE FA0F    ANI #0F Get LSD
C0 F4      ADD D+M(RX)
C1 30AE    BR DECH-1 Exit
C3 XX      .. Hexadecimal to Decimal Subroutine
C4 9F      GHI RF Get RF.1 into D
C5 D4      SEP R4 Return to main
C6 BF      HEXD:PHI RF Enter Byte
C7 F800    LD1 00 Initialize
C9 AB      PLO RB
CA AF      PLO RF
CB 9F      GHI RF Bring back byte
CC FF64    SMI 100 Subtract 100
CE 3BD3    BNF *-5 If less than 0
D0 1F      INC RF
D1 30CC    BR *-5 Go back to Subtract 100
D3 FC64    ADI 100 Add 100 to get LSD
D5 FF0A    SMI 10 Subtract 10
D7 3BDC    BNF *-5 If less than 0
D9 1B      INC RB
DA 30D5    BR *-5 Go back to subtract 10
DC FC0A    ADI 10
DE BF      DL: PHI RF
DF 8B      GLO RB
E0 32C4    BZ HEXD-2 Exit
E2 9F      GHI RF
E3 FC10    ADI #10 Add hex 10
E5 2B      DEC RB
E6 30DE    BR DL Loop until D=0
E8 xx      Temporary storage
.. xx
FF xx

```

### Square Roots

The mathematical capabilities of Tiny Basic are really rather limited: no fractions, no trig functions, not even a square root. Who needs them? Maybe you do. So this is a short program to calculate the square root of a number.

If you write out the perfect squares in a row like this

0 1 4 9 16 25 36 49

you can subtract each one from the next square in the series:

1 3 5 7 9 11 13

Notice that the differences are simply the odd numbers in sequence. In other words, if I start with zero and successively add the odd numbers, 1, 3, 5, 7, and so on, I will get all the perfect squares.

We can use this fact to compute the square root of a (small)

number, by subtracting successive odd integers from it until it goes below zero. Here is a Tiny Basic program to do this:

```
100 REM COMPUTE SQUARE ROOT
110 INPUT N
120 IF N<0 THEN GOTO 110
130 PRINT "THE SQUARE ROOT OF ";N;" IS ";
140 F=1
150 IF N<327 THEN F=10
160 N=N*F*F
190 I=-1
200 REM MAIN LOOP
210 I=I+1
220 N=N-I-I-1
230 IF N>=0 THEN GOTO 210
300 REM PRINT THE RESULT
310 PRINT I/F;
320 IF F>1 THEN PRINT ".*";I-I/F*F;
330 PRINT
340 GOTO 110
```

This program should fit into the minimum Tiny Basic configuration (see Page 2, this issue), if you take out the REM lines.

Another way to calculate a square root is by successive approximation, a kind of "cut and try" technique known as

Newton's Method. In this you make a first guess at a root, then divide it into the original number and average the divisor and quotient for the second guess. This is repeated until you cannot get any closer. Newton's Method is best when the numbers you are working with are large, or if the computer has a good, fast, hardware divide.

We can modify the program to use Newton's method:

```
190 I=1
210 R=I
220 I=(N/I+I)/2
230 IF I<>R THEN GOTO 210
```

Notice that this version only changes the inner loop. For some numbers this program never finishes. Can you figure out why? In case you cannot, I'll give the answer next month.

You may wonder what the purpose of F is. If the number is small enough, we can get a more precise result by a process known as "scaling". That is, multiplying the number by a scale factor, then dividing the result by (in this case) the square root of that scale factor. The scale factor is 100, and if 100 times the number is still in the Tiny Basic range (i.e. is less than 32768) then we can get one more decimal place of accuracy by scaling.



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